Resilience or Robustness: Identifying Topological Vulnerabilities in Rail Networks

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Abstract

Many critical infrastructure systems have network structure and are under stress. Despite their national importance, the complexity of large-scale transport networks means we do not fully understand their vulnerabilities to cascade failures. The research in this paper examines the interdependent rail networks in Greater London and surrounding commuter area. We focus on the morning commuter hours, where the system is under the most demand stress. There is increasing evidence that the topological shape of the network plays an important role in dynamic cascades. Here, we examine whether the different topological measures of resilience (stability) or robustness (failure) are more appropriate for understanding poor railway performance. The results show that resilience and not robustness has a strong correlation to the consumer experience statistics. Our results are a way of describing the complexity of cascade dynamics on networks without the involvement of detailed agent-based-models, showing that cascade effects are more responsible for poor performance than failures. The network science analysis hints at pathways towards making the network structure more resilient by reducing feedback loops.

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1 Introduction

Cascade delays and cancellations on rail transport can cause devastating economic damage and dent consumer satisfaction. Existing knowledge focuses either on improving operational practices or considers pure topological analysis. We advance this, by considering both real passenger travel flows and the network topology together. This creates a stronger understanding of its dynamic vulnerability and resilience. In earlier years, research largely focused on improving specific functionalities in rail systems; and more recent research has focused on the relationship between the general network topology and whether this has macroscopic bearing on the overall system performance [1].

The efficiency of transport networks has been related with their resilience [2] and the different types of topologies have been analysed, comparing the network geometry and the level of connectivity. However, these studies predominantly focus on the pure topological characteristics of a graph (Guo and Cai [3] and Cañizares et al. [4]).

1.1 Related Work

1.1.1 UK rail network

The UK rail network transports more than 1.7 billion passengers per year, of which 1.1 billion passengers commute in the London area [33]. According to the Office of Rail and Road [34], in the last year in the London area, only 86.9% of passenger trains arrived on time and 4.8% of the journeys experienced cancellations or significantly lateness. Often these delays are interrelated and the relationship between cascade effects and network dynamics is not well understood.

In current literature, most of the proposed studies consider natural or man-made disasters, but they do not consider the stress of the network during the peak-hours and how the structure of the network created by the massive flows of people can influence their ability to maintain a good service. For example, several graph-based approaches have been proposed to improve the performances by revising the design and maintenance of the rail networks [36, 39], but do not consider dynamic passenger flows. Other studies focus on specific extreme scenarios [37, 39] or unfavorable conditions [35] that cause disruptions.

As our data shows, under the same external conditions, the major rail companies in and around London reach dramatically different performance levels. In this work we hypothesize that this difference can in part be attributed to the peak passenger demand. A coupling relationship between flow...
and network structure can tease out the indicative measures that correlate strongly with overall performance.

1.1.2 Vulnerability of Transport Networks

The concept of vulnerability of transportation network, introduced in the literature by Berdica [40], is generally defined as the susceptibility to disruptions that could cause considerable reductions in network service or the ability to use a particular network link or route at a given time. Many have applied general network science disruption analysis. For example, several studies [43, 44, 45] have been conducted for modelling railway vulnerability with promising predictive results. Bababeik et al. [42] recently proposed a mathematical programming model that is able to identify critical links with consideration of supply and demand interactions under different disruption scenarios. Recent work has also used graph properties to infer interaction strengths and use an epidemic spreading model to predict delays in railway networks [54].

1.2 Innovation

In this paper, we take a systems-of-systems approach by applying complex network analysis to transport networks. Unlike prior studies that focus only on the topological aspects of the network, we consider several important additional aspects which attempt to match our analysis to reality. First, we consider passenger volumes during morning commuter or rush-hour, which weights the network and adds directionality. Considering the morning rush-hour is important because most of the delays and the highest economic impact of delays occur during these times. Second, we consider a multiplex of different urban overground, regional, and national rail services (both together and separately). As a result, we have a weighted and directed multiplex network, which requires more sophisticated network analysis methods to uncover its resilience and robustness to cascade failures. Finally, we map our network resilience and robustness results to actual railway performance figures of delay and cancellation statistics and consumer satisfaction.

1.3 Analysis

Vulnerability is a major problem in the study of complex networks and it can be regarded as the susceptibility of a networked system to suffer important changes in its structure and dynamic functions under any form of disruption. When such disruptions affect the internal state of the nodes (e.g., stations)
or links (e.g., train lines) of the network, it becomes important to predict the extent of such perturbations under the perspective of dynamical systems (e.g., linear stability analysis); throughout this paper, we refer to this problem as the study of resilience. Resilience is important for understanding cascade effects that suppress the performance of the network, such as cascade delays due to signal failures or poor scheduling. In plain terms, resilience describes the problem of a train from A to B that is late, which will affect the ensuing service B back to A using the same train. On the other hand, when the perturbations involve some sort of attack or out-right failure (e.g., a disruption in a station due to someone walking on to the tracks or a signal failure), the challenge tends to be in studying the resulting connectivity loss and secondary loss of functionality in neighbouring stations. We refer to this as the robustness problem, which is very different from the aforementioned resilience. In plain terms, robustness considers when a train from A to B will be halted if the track in between is blocked or station B is closed.

The concepts of resilience and robustness on networks admit various interpretations and definitions [5, 6]. A generally accepted definition of stability is for when the system performance returns to a desirable state. For homogeneous linear stability, one might equate resilience with equilibrium points and look at the leading eigenvalue of the Jacobian matrix [7]. When linear stability is not suitable due to complex dynamics, many authors [47, 48, 50, 49, 51] have studied system resilience from different perspectives. Some consider the dynamic response (e.g., time to recovery) of the whole system after a specific disruption [51], whilst others use random perturbations to numerically quantify system response [52].

In terms of robustness, a common definition is the number of nodes that must be removed in order for the network to break down is a popular measure of its robustness [8]. Whilst such approaches depend strongly on assumptions about the system, it generally maps well to railway systems [46].

However, such approaches depend strongly on assumptions about the system, such as details of the dynamics or the number of neighbours required for a node to function. In this work we make use instead of recent advances in ecological system analysis to study resilience and robustness, which can be obtained directly from the adjacency matrix (even for weighted and directed networks) and have been found to be good proxies for resilience and robustness in ecosystems (Fig. 1). While there are certainly differences between ecosystems and rail systems, both are essentially transport networks in which either biomass or passengers flow from sources (plants, or home towns) through various intermediary nodes, and end in sinks (top predators, or work places).
Figure 1: We reconstruct the major rail networks under stress conditions considering the morning journeys (a) and we measure the topological characteristics of these networks: the uninteresting flows are removed (b), then the resilience (c) and robustness (d) of these networks are analysed.
1.3.1 Resilience of Weighted and Directed Networks

We introduce a parameter for quantifying the resilience of weighted directed networks measuring their *trophic coherence*. Trophic coherence is a property of directed graphs that defines how much a graph is hierarchically structured. The rationale is that hierarchical systems have fewer feedback loops and are less likely to suffer from cascade effects.

When networks are modeled as a discrete linear time invariant (LTI) system with a defined input and output [53], the dynamic response stability is defined by the location of roots of its transfer function (negative domain). In such a case, absence of feedback loops ensures stability. Presence of feedback loops will cause non-zero roots and risk instability. When we consider a complex network with $\sim N^2$ input output combinations, the transfer function cannot be defined. As such, we measure the overall network incoherence, which is a compressed figure of merit for how many feedback loops exist [9, 10]. Johnson et al. [9, 10] proved that "a maximally coherent network with constant interaction strengths will always be linearly stable", and that it is a better statistical predictor of linear stability than size or complexity. We measured the coherence of the network through the *incoherence parameter*, a measure of how tightly the trophic distance associated with edges is concentrated around its mean value (which is always 1) [9].

In order to define trophic coherence in a directed network, the first step is to define basal nodes (i.e., nodes that predominantly supply energy – high out-degree and low in-degree). That is to say, stations with a high trophic level receive passengers while stations with a low trophic level provide passengers. Thus, basal nodes are likely to be home train stations of commuters.

Defining trophic coherence in real data networks requires some pre-processing: unlike the already studied networks in other works (e.g., food webs [11, 12]), the London urban rail network in peak-hours does not have predefined basal nodes (i.e., nodes with in-degree of 0). In transportation, this means that there is always a non-zero passenger counter-flow traveling from urban to the countryside stations during the morning rush hour. In order to distill the basal nodes from the data, we developed and tested two different approaches (we show these techniques in Section 4, page 15) to identify the basal nodes in networks where they do not naturally exist. In other words, many weighted networks do not have apparent energy sources. In the first proposed approach, we apply **basal node enforcement**, whereby basal nodes are selected based on their network centrality characteristics (e.g., out-degree). The trophic level of the remaining nodes is then computed using the standard formula (equation 2). In the second proposed approach, we apply **passenger flow filtering**, which identifies the basal nodes by reducing the
connectivity of the network. We sequentially reduce the connectivity (by increasing a flow threshold) until basal nodes emerge naturally.

1.3.2 Robustness

The objective in this case is to use both proxy and direct measures of robustness. Direct measures are random or targetted node removal. However, as robustness is not rigorously defined, proxy measures may yield a more holistic insight: a variety of robustness measures are used to compare against real railway performance. As such, we are using a variety of robustness measures to establish a wider evidence base.

Firstly, as a proxy, we evaluate its core and periphery meso-scale structure. The core periphery ratio gives a scalable and compressed understanding of robustness, and the argument is formalised by Borgatti et al. [13]. Another proxy measure is the rich-club coefficient [14, 15, 11, 12]. Secondly, we evaluated the robustness directly by performing sequential node removal [16]. The nodes of the rail networks are removed randomly and the network connectivity is then studied evaluating the size of the largest strongly connected component.

1.4 Data

In this study we analysed the rail network under demand stress conditions (morning rush-hour). The commuter paths are computed considering the information relative to places where people live and work provided by the UK National Census Transformation Programme [17]. The optimal travel paths were provided by the National Rail (including rail services through underground tunnels, but not including the underground/subway system) through their TransportApi service [18]. Given an origin and a destination stations, the TransportApi service provides all the information about the travel, including the intermediate stop stations. We first check if rail travel is required for a person to go from home to work, and if so, we compute their optimal journey and use this data to weight the network. In the current study, only the travels that start and end in a bounding area of 80 km from central London have been taken into account (this approximately covers Cambridge to the north, Oxford to the northwest, Reading to the west, and Brighton to the south). It roughly represents all 1 hour commuter paths, which is the national standard according to ONS [19].

The resulting dataset represents the flows of people in morning peak-hours on the rail network (will be made available on Dryad), when they travel from their homes to their places of work. Each journey is defined
as a set of two or more stations (in case of intermediate stops of the train all the intermediate station are included). The dataset is transformed in a directed weighted graph where the nodes are the train stations, the edges are the weighted flows of passengers and a journey is an ordered set of nodes that includes the departure station, the arrival station and any intermediate station (if the train stops, as we consider the service class of the train).

A directed graph is defined \( G = (N, E) \) as an ordered pair where \( N \) is a set of nodes (i.e., stations) and \( E \) is a set of ordered pairs of nodes, called edges (i.e., trains that go from a node to the following). When, in our graph, one or more passengers are going from node \( i \) to node \( j \) (or these two nodes are intermediate stations of the travel), an edge \( e_{ij} \) is added to the graph. The weight of this edge is the sum of all the passengers of the journeys that include travels from node \( i \) to node \( j \). The directed graph of the passenger flows during a morning peak-hours is shown in Figure 2. We show the whole multiplexed network, as well as some examples of the individual sub-networks comprised of urban overground (London Overground), regional links (Thameslink), and national services (e.g., Southern rail).

Figure 2: Directed graph of passenger flows during a morning peak-hours.
2 Results

The delays in a rail network and, more in general, the performances of the service are influenced by the topological structure of the network. The notion is that a more resilient and/or robust network should guarantee lower cascade delays and faster recovery in case of disruptions. In this chapter the different multi-scale rail networks (this includes both local London overground rail services and national major rail services) in a morning peak-hours are analysed separately and the results are compared with the Public Performance Measures provided by the ORR (Office of Rail and Road) \[21\], an independent regulator that monitor the rail industry’s health and safety performance. ORR holds Network Rail \[22\] the company that, with 20,000 miles of track, owns, operates and develops Britain’s railway infrastructure.

In particular, 2 performance measures are used in our comparison:

PPM: the Public Performance Measure combines figures for punctuality and reliability into a single performance measure. Usually, it shows the percentage of trains which arrive at their terminating station within 5 minutes (for London and South East and regional services) or 10 minutes (for long distance services) \[23\]. In this paper, for the sake of clarity, we define the oPPM with the opposite value of the PPM ($oPPM = 100\% - PPM$). oPPM is the percentage of trains which not arrive at their terminating station within 5 or 10 minutes (depending on the distance).

CaSL: the Cancellation and Significant Lateness is a percentage measure of scheduled passenger trains which are either cancelled (including those cancelled en route) or arrive at their scheduled destination more than 30 minutes late \[23\].

The performance measures utilised in this paper are referred to the year 2017 (key statistics by train operating company (TOC) - 2016-17 \[24\]). To provide statistically significant results (small networks are more sensitive to local functional effects than macroscopic topological structure), we considered the 5 companies with the highest number of nodes in the network, excluding companies with very simple network structures (e.g., Heathrow Express has only one line). The companies taken into account and the number of stations are shown in Table 1.

4 out of 5 analysed rail companies (see Figure 3) show a strong correlation between these 2 measures while in one case (Great Western Railway) these values are not correlated, possibly meaning that this company often has little
delays (low resilience) but generally does not have major disruptions (high robustness).

The Pearson Correlation Coefficient (PCC) \([25]\) is used to establish if there is a correlation between the topology parameters of the network and the performance measures. PCC has a value between +1 and -1, where 1 is total positive linear correlation, 0 is no linear correlation, and -1 is total negative linear correlation. Two variables with a correlation coefficient greater than 0.7 are considered highly correlated, while they are considered moderately correlated when the PCC coefficient is between 0.3 and 0.7.

### 2.1 Trophic Incoherence Analysis

The degree to which the rail networks are coherent (or incoherent) can be investigated by comparison with a null model. We use the basal ensemble expectation \(\tilde{q}\) as a null model to compare the incoherence parameter of our rail networks. The trophic incoherence measure \(q/\tilde{q}\) has a value close to 1 when a network has a trophic coherence similar to a random expectation, it has a value lower than 1 when the network is coherent while it has a value greater than 1 when the network is incoherent (more details and the computation of the basal ensemble expectation are provided in Section 4.4.1.4.1.1).

The incoherence coefficient \(q\) of a morning peak-hours network is computed using the passenger flow filter method with different flow filtering thresholds, between 1 and 4, with a granularity of 0.5 (details on the selection of the methodology and the parameters are given in Section 4.4.1). The average of the computed incoherence parameters is compared with the relative oPPM and the CaSL measures and shown in the following figures along with their standard deviation.
Our results exhibit a highly positive correlation between the trophic incoherence of the network and the Public Performance Measure (PCC = 0.98), suggesting that there is a high correlation between the resilience of a rail network and the probability of its trains to arrive at their terminating station on time. There is also a high positive correlation between the trophic incoherence and the Cancellation and Significant Lateness measure (PCC = 0.92), evidencing also a correlation between low resilience and the percentage of trains either cancelled or that arrive to their destination with more than 30 minutes late. The trophic incoherence measure \( \frac{q}{\tilde{q}} \) compared with oPPM and CasL is shown in details in Figure 4: more coherent networks (low \( \frac{q}{\tilde{q}} \)) are generally associated with lower delays (oPPM) and cancellations (CaSL).

Figure 4: oPPM and CaSL compared with trophic incoherence parameter of each rail company.

Figure 5: oPPM and CaSL compared with Rich-Club coefficient of each rail company.

2.2 Rich-Club Coefficient Analysis

The rich-club coefficient of a network is an indicator of the robustness of a network, in particular we say that networks with a rich-club coefficient greater than 1 are characterised by the rich-club phenomenon. The rich-club coefficient measures how many nodes of degree at least \( k \) are connected in a graph (how many edges are present), normalized by the maximum number of possible connections between them (maximum number of edges) in a complete graph (details are provided in Section 4.4.2.4.2.1). In this paper we used the standard definition of rich-core, classifying the nodes according to their degree.

We compared the highest coefficient reached (considering all the possible \( k \) values) for each company with its performances metrics (Figure 5). Our results shows that, even if there is a moderate correlation between the value of the rich-club coefficient and the performances (PPM has PCC = 0.62 and
CaSL has PCC = 0.55), there is no evidence of any correlation between the presence of the rich-club phenomenon and the service performances of the companies.

2.3 Core Size Analysis

The size of the core of a network compared to its periphery represents the percentage number of well connected core stations, versus the sparse periphery stations (as we discussed, intuitively a network with a bigger core has more connections between the stations and, thus, a higher robustness to disruptions). In this section we compare the percentage of core nodes of each company network, computed using degree and trophic level for ranking, and its oPPM and CaSL measures.

Our findings show that there is a moderate positive correlation between the size of the degree core (PCC = 0.38) and the trophic core (PCC = 0.59) of a company network and the oPPM. Instead, there is no correlation with the CaSL (degree core PCC = -0.09, trophic core PCC = 0.28). A comparison between the degree and trophic core of the companies and the performance measures is shown in Figure 6 and Figure 7.

Figure 6: oPPM and CaSL compared with degree core of each rail company.

Figure 7: oPPM and CaSL compared with the trophic core of each rail company.

2.4 Removal of Random Nodes Analysis

We attacked the network removing nodes and analysing the size of the largest component. The experiments were repeated several times removing the nodes randomly: in Figure 8 is shown the average size of the largest component of each company and its standard deviation. The threshold value, \( T = 50\% \) (dashed red line in the Figure), has been chosen as the limit value to consider a network 'alive', the percentage of nodes required to destroy a network
(connectivity ≥ 50%) is defined as its robustness to attacks. The robustness to attacks is compared with the performance measures of the companies (Figure 9). Our results show a strong correlation between the robustness to attacks and the CaSL measure (PCC = 0.83) and a moderate correlation (PCC = 0.58) with the oPPM measure.

The correlation analysis has been extended to other significant company statistics (Figure 10) showing that oPPM and CaSL are correlated with some of the other values, but the incoherence ratio $q/\bar{q}$ is indeed one of the most significant. The robustness to attacks has shown to be a good indicator for cancellations and significant delays (CaSL). The size of the core (both degree and trophic cores) and the rich-club phenomenon do not provide any significant correlation with performances. The size of the rail network in terms of the number of employees and stations also has a strong correlation to oPPM, which is likely to be the effect that larger networks are more likely to have feedback loops and incur cascade effects.

3 Discussion

In this work we proposed a study of the London’s urban rail network under stress conditions, during the morning peak-hours. We represented the major companies rail networks as weighted directed graphs, where the nodes indicates the stations, the edges indicates the flows of people and the weights of the edges the number of people traveling on that segment. If two stations
Figure 10: **Pearson Correlation Coefficient** between different measures and indicators. Incoherence parameter / basal ensemble ($q/\tilde{q}$), incoherence parameter ($q$), size of degree core (Size Degree-Core), size of trophic core (Size Trophic-Core), rich-club phenomenon (Rich-club), Robustness to attacks (Attacks), 100% - Public Performance Measure (oPPM), Cancellations and Significant Lateness (CaSL), number of employees, number of stations, number of trains, number of passengers and multiplex stations (stations that connect different companies).
are connected but there are no passengers travelling in the morning, these stations are considered disconnected (there is not an edge between these nodes).

We studied the resilience and the robustness of these networks drawing inspiration from techniques used to study natural complex networks, such as food webs. Our results suggest that the resilience indicators \((q, \frac{q}{\bar{q}})\) have a strong correlation with the performance parameters PPM (Public Performance Measure) and CaSL (Cancellations and Significant Lateness). Conversely, the different robustness indicators (size of the core and rich-club phenomenon) are not significantly correlated with these measures, although the robustness to attacks is correlated with the CaSL measure.

There are a number of further interesting research that the community can action using our data and building on our methods. Firstly, an interesting improvement in the current work would be the introduction of a theoretical model that could help assessing the role of noise – understood as deviations from shortest path routing protocols between origin and destination [26] – in the design of resilient flows in complex networks. In particular, we hypothesize the existence of trade-off between resilience and travel distance mediated by the amount of noise present in the network flows: strictly shortest path protocols will tend to overload links when the network is attacked [27, 28], an effect that could be mitigated by adding randomness in the rerouting. Intuitively, artificially inducing random behaviour (noise or dynamic routing with randomness) in human behaviour may even alleviate overall congestion and thus achieve a lower global travel time. Secondly, we have in this paper compressed complex network dynamics into a simple coherence metric. Given we want to maximize coherence, the reverse open question is how best to make the minimum number of changes in scheduling and/or graph to maximize a positive step change in coherence. This problem is open ended, because there are clearly local economic reasons for loops in transport and only a limited number of train carriages to serve them.

4 Methods

4.1 Trophic Coherence and Resilience

The trophic level of a node \(i\), called \(s_i\), is defined as the average trophic level of its in-neighbours, plus 1:

\[
s_i = 1 + \frac{1}{k_i^{\text{in}}} \sum_j a_{ij} s_j,
\]  

(1)
where $a_{ij}$ is the adjacency matrix of the graph and $k_i^{in} = \sum_j a_{ij}$ is the number of in-neighbours (in degree) of the node $i$. Basal nodes $k_i^{in} = 0$ have trophic level $s_i = 1$ by convention.

By solving the system of equations (1), it is always possible to assign a unique trophic level to each node as long as there is at least one basal node, and every node is on a directed path which includes a basal node [9]. In our study the trophic level of a station is the average level of all the stations from which it receives passengers plus one. For this reason, stations near residential areas in the suburbs will have lower trophic level than those close to business areas in the centre.

Each edge has an associated trophic difference: $x_{ij} = s_i - s_j$. The distribution of trophic differences, $p(x)$, always has mean 1, and a network will be more trophically coherent the smaller the variance of this distribution. We can measure trophic coherence with the incoherence parameter $q$, which is simply the standard deviation of $p(x)$:

$$q = \sqrt{\frac{1}{L} \sum_{ij} a_{ij} x_{ij}^2 - 1} \tag{2}$$

where:
$L = \sum_{ij} a_{ij}$ is the number of connections (edges) between the stations (nodes) in the network. A perfectly coherent network will have $q = 0$, while a $q$ greater than 0 indicates less coherent networks.

In the morning peak-hours rail networks object of our study there are not natural basal nodes. As they are a requirement to solve the equations for the trophic levels computation, we defined two methodologies to identify them: the basal nodes enforcement and the flows filtering.

### 4.1.1 Comparison with null model

The degree to which empirical networks are trophically coherent (or incoherent) can be investigated by comparison with a null model. The basal ensemble expectation $\bar{q}$ can be considered a good approximation for finite random networks [10]. We use this parameter as a null model to compare the incoherence parameter of our empirical networks.

The basal ensemble expectation for the incoherence parameter is [10]:

$$\bar{q} = \sqrt{\frac{L}{L_B} - 1} \tag{3}$$

where:
$L = $ number of edges in the network.
$L_b =$ number of edges connected to basal nodes.

The ratio $q/\tilde{q}$ is used to analyse the coherence of the network: a value close to 1 shows a network with a trophic coherence similar to a random expectation. Values lower than 1 reveal coherent networks, while values greater than 1 incoherent ones.

Johnson and Jones [10] found that food webs are significantly coherent ($q/\tilde{q} = 0.44 \pm 0.17$), metabolic networks are significantly incoherent ($q/\tilde{q} = 1.81 \pm 0.11$) and gene regulatory networks are close to the random expectation ($q/\tilde{q} = 0.99 \pm 0.05$).

4.1.2 Basal Nodes Enforcement

The first technique used to select the basal nodes revolves around the enforcement of the desired number of basal nodes, selecting them according to some properties of the nodes. This technique enforces a predefined number $EN$ of nodes to be basal nodes (their trophic level is imposed 1). The nodes to be enforced are selected according to their similarity to real basal nodes, namely the nodes with the lowest ratio between incoming and outgoing edges. More formally, the $k^{\text{out}}/k^{\text{in}}$ ratio is computed for all the nodes, then the trophic level of the $EN$ nodes with the lower ratio is enforced to 1 ($s_i = 1$). If parts of the network are not connected to basal nodes, only the largest strongly connected component considered. This technique maintains the structure of the network intact (it does not add/remove nodes or edges) but, instead, it does not take into account its natural topology when selecting the basal nodes, making the selection artificial: the selection of the number of basal nodes is artificially defined by the user and does not evaluate the ideal natural number of basal nodes present in the network.

4.1.3 Flows Filtering

In the analysis of the morning peak-hours commutation, the factors that determine the stability of the network depend on the major flows of people (from home to work commutation). The paths with just a small portion of commuters can thus be ignored. To remove these paths, a threshold $T$ for the detection of the major flows is defined: when two nodes $i$ and $j$ are connected with two edges $a_{ij}$ and $a_{ji}$, the edges whose ratio $a_{ij}/a_{ji}$ is below the threshold $T$ are deleted.

With this approach it is possible to remove those loops in the network that are not relevant for the peak-hours analysis (e.g., 100 people going from node $i$ to $j$ and only 1 going from $j$ to $i$, the edge $(j,i)$ can be removed.
without degrading the quality of the peak-hours flows study). If $T \geq 1$, for each pair of nodes, only the edge with the highest weight is maintained, and only if it is greater than the other (otherwise both the edges are removed). The larger $T$, the more the direction of the flow from one station another has to be predominant compared to the reverse. If $T < 1$, both edges could be possibly maintained if their flows were balanced (the lower $T$ the more unbalanced the flows).

With this technique the basal nodes are not enforced but rather naturally emerge from the change in the structure of the network (i.e., the edges with a low impact on the study are removed from the network). The higher the threshold, the more edges are removed. However, the threshold has to be accurately chosen because a high threshold could lead to the removal of interesting flows, reducing the information in the graph and providing incomplete results.

4.2 Core-Periphery and Robustness

The study of the core-periphery structure of the network is used to identify the densely-connected stations where people can choose more than a path to reach the destination in contrast to sparsely-connected stations which can cause a major interruption of the service in case of disruptions. The core of a network [29] is computed ranking all the nodes in a network and then counting the number of connections (edges) they have with higher ranked nodes. The node with the highest number of high-level connections is the core border. All the nodes with a higher ranking than the border node along with the border node itself compose the core of the network, the other nodes are the periphery. A big core suggests several different ways to reach the majority of the nodes and accordingly a more robust network.

4.2.1 The Rich-Club Phenomenon

To study the robustness of the networks we analysed the rich-club phenomenon [30]. It is characterized when nodes of higher degree are more interconnected than nodes with lower degree. The presence of this phenomenon may indicate several interesting high-level network properties, such as its robustness. More precisely, this behaviour appears when nodes larger than $k$ are more densely connected among themselves than the nodes with degree smaller than $k$ [31]. This is quantified by computing the rich-club coefficient across a range of $k$ values, if this value is greater than 1 for some $k$ the network is characterized by the rich-club phenomenon.
The rich-club coefficient for a given network $N$ is usually defined for the degree of the nodes, but it can be generalized to other metrics (e.g., the trophic level). We converted the morning peak-hours directed graph to an undirected one to be consistent with the standard rich-club definition. The generalized formula to compute the rich-club coefficient is:

$$\phi(r) = \frac{2E_{>r}}{N_{>r}(N_{>r} - 1)},$$  

(4)

where only nodes with a richness measure (e.g., node degree rank) of at least $r$ are considered, $E$ is the number of edges and $N$ the number of nodes.

### 4.3 Data-driven Analysis of the Methods and the Parameters

#### 4.3.1 Resilience and Trophic Coherence

In this Section we discuss the properties of the basal nodes selection methods. Our empirical results show that, on one hand, the node enforcement always produce incoherent networks (even with a large number of stations enforced we have a high $q/\tilde{q}$ ratio ($> 2$)). On the other hand, the passenger filtering technique can achieve a stable network and low incoherence level: eliminating links with a passenger in - passenger out ratio greater than 3.5 creates a network with a $q/\tilde{q}$ of 0.6 (0.4 with passenger in - passenger out ratio $> 8$). The latter method is therefore selected, not only because it makes intuitive sense (e.g., a small number of counter-flow passengers is regarded as noise that is filtered out), but also because it creates a network with clear trophic levels that match our qualitative knowledge of how passengers travel.

It is crucial for a reliable analysis to select filter values that are reasonably representative of some underlying data structure. Here, we look for the minimum filter value (so we do not remove too much data), such that the measure of interest (e.g., coherence or core size) is invariant to further filter value changes. Referring to Figure 11, passenger flow filter method is reasonably stable compared to enforced basal node method. Whereas, for trophic core size, we can see that a relatively larger value of passenger flow filtering is required. For degree core, we can see that any reasonable filtering or basal node enforcement produces a reliable answer.

Like for the overall rail network infrastructure, the rail network in morning peak-hours of each separated company may not have basal nodes, thus to compute its trophic incoherence we firstly identified them using one of the techniques described. The flows filtering methodology, that has proved to
be a better technique for our work, is used to identify the basal nodes and compute the incoherence parameter $q$ for each company network. To obtain more homogeneous results the flow filter is applied to all the networks, even if some basal nodes are already existing.

According to the previous results, flow filtering threshold $T$ between 1 and 4 provides the best conditions to study the network, filtering the edges that do not represent the studied behaviour (small counter-flows) without modifying significantly the network structure (higher thresholds may also remove interesting flows).

4.3.2 Robustness and Core-Periphery

![Graphs showing incoherence, trophic core-periphery ratio, and degree core-periphery ratio.]

Figure 11: Behaviour of the three major resilience and robustness measures used in this work. The basal nodes are selected with several parameter values of the two techniques proposed: the basal nodes enforcement parameter is the number of nodes (red) and the flows filtering parameter is the filter threshold.

The study of the core-periphery structure of the network is used to identify the densely-connected stations where people can choose more than one path to reach the destination in contrast to sparsely-connected stations which can cause a major interruption of the service in case of disruptions.

We investigate if a higher percentage of nodes belonging to the core is related to better performances.

Two metrics are used to evaluate the core-periphery of the network. The first metric is relative to the classic definition of network core and rich-club phenomenon: the node degree. The second metric is specific of the studied rail network: as shown previously, in the morning peak-hours rail network, most people travel from the low trophic level loosely connected periphery stations to the high trophic level well-connected core stations. For this reason, in order to identify the core of the morning peak-hours network, the second
metric chosen to rank the stations (nodes) is their trophic level, computed using both the basal node enforcement technique and the passenger flow filter technique.

In this section we provide a first analysis of the robustness of the morning peak-hours network counting the percentage of nodes that belong to the core of the network, using the degree and trophic level metrics. We will refer to the **degree core** as the core of the network computed ranking the nodes according to their degree, we will refer to the **trophic core** as the core of the network computed ranking the nodes according to their trophic level.

A comparison of the two approaches is shown in Figure 11(b) and 11(c). Our results show that the trophic core is much bigger than the degree core. The former core is always between 20% and 90%, the latter is around 5% of the overall nodes. Generally, the passenger flow filter technique generates a network with a bigger core than the enforced basal nodes one. Our experimental results suggest that, while the network is generally not well connected with the higher degree stations (the centre of the network), it is well connected with the stations with a higher trophic level (the morning destination of the passengers).

### 4.3.3 Which Basal Nodes Identification Method?

In the London Urban Rail network during the morning peak-hours there are no natural basal nodes, so we provided two techniques to artificially select them. Generally speaking, the **Enforcing basal nodes technique** does not modify the structure of the network and can be a good technique when all the edges are important or the graph is unweighted. Contrariwise, in scenarios where the difference of the edge weights is significant or the focus is on a certain kinds of network behaviour, it may be worth using other approaches. In our study, where the focus is on the morning peak-hours passenger flows, the more reliable approach consists in using the **Flow filtering** method, with a threshold on the *passenger flows* that removes the small counter-flow (e.g., people that live in the centre and work in the suburbs) in order to evidence the mass commutation that causes the major stress on the network. With the latter approach and a fair passenger filter threshold we obtain a network with a trophic coherence similar to a null model. This shows that the rail network during morning peak-hours has a coherence similar to the random expectation, not as incoherent as suspected by enforcing the basal nodes (which does not remove the irrelevant edges).
5 Appendix

5.1 Alternative Data Testing

The rail performance statistics used in the main paper sections are relative to the year 2016/17, and here we show this is not a spurious result by using alternative years data. We perform the same analyses using the rail statistics relative to two other years. On one hand, we compare the resilience and robustness measures with the statistics relative to the year 2012 (2011/12 Q4), the same year of the census dataset used to create the morning peak hours network. On the other hand, we compare the same measures with the rail statistics relative to the year 2018 (2018/19 Q1), the latest currently available. All the performance statistics are publicly available on the Office of Rail and Road website [21].

5.1.1 Year 2011/12

In this section we discuss the correlation between the resilience and robustness measures and the PPM and CaSL relative to the year of the census used to create the morning peak hours network (2011/12 Q4).

The new results (Figure 12) confirm the ones in the previous analysis: as for 2017, we found a correlation between the normalised trophic incoherence ($q/\tilde{q}$) and the oPPM and CaSL measures (Figure 13).

5.1.2 Year 2018/19

In this section we compare the resilience and robustness measures of the different rail networks with the latest statistics available (2018/19 Q1). We also extended the study to two more rail companies: the Thameslink and the Southern Railway. Moreover, South West Train has been renamed as South Western Railway. The new results (Figure 14) confirm the ones in the previous analysis: as for 2017, we found a correlation between the normalised trophic incoherence ($q/\tilde{q}$) and the oPPM and CaSL measures (Figure 15).

5.2 Case study: How could adding or removing a single link improve the rail network performances?

In this section we propose a case study on the Thameslink rail network morning peak hours. Aim of this study is to show how a single modification in the network structure could improve the performance of the network. Given
the correlation between the trophic incoherence and the rail performance measures (oPPM and CaSL) evidenced in this paper, we propose possible network modifications that reduce the first in order to improve the seconds.
We look for changes in the network that improve the trophic coherence without causing major disruptions in the service. In this case study, we focus on a single modification: we add and remove a link and we measure the resulting trophic incoherence, we then analyse in details a scenario with a link removed or one with a link added.

5.2.1 Removing a link

Removing a link in the morning peak hours network can lead to major disruption (e.g., isolate parts of the network) and lengthen the commutation time. We detect these possible problems by measuring two parameters:

- Closeness: Closeness measure the centrality of a node in a network and it is thus related to the commutation time. Closeness:

  \[ C_i = \frac{n - 1}{\sum_j d_{i,j}} \]

  We compute the average of the nodes’ closeness in order to have a standard measure for comparison. We use the average closeness (even if the majority of the nodes are not affected by a single modification), because we are interested in how the whole network is affected by a single link removal (the bigger is the decrease the more the average commutation time increases).

- Largest connected component: We want to avoid isolation of parts of the network, for this reason we measure the size of the largest connected component: if it decreases, the possible link removal is not taken into account.

We removed each node one by one, each time computing the new trophic incoherence, the average closeness and the size of the largest connected component and we selected the link that provides the best improvements without breaking the discussed constraints.

Our analysis suggests that removing the link from Denmark Hill to Elephant and Castle provides the best incoherence gain without isolating parts of the network and without a huge increase of the commutation time.

This modification improves the trophic incoherence by 41.4%. It reduces the average node closeness of around 10%.
Figure 16: Stations trophic level. On the left the trophic level of the Thameslink stations, on the right the trophic level of the same stations when the link from Denmark Hill to Elephant and Castle is removed.

Trophic Incoherence Analysis In Figure 16 we show the incoherence gain of the Thameslink network when the Denmark Hill - Elephant and Castle link is removed. Denmark Hill station, the node with the highest trophic level, decreases it from 22.8 to 19.4. While this reduction influences the trophic incoherence, the main contribution in its abatement is due to another factor: loops are linked to trophic coherence and removing them improves the network trophic structure. In our case study, we detected a loop in the railway section between Blackfriars and Bromley South (see Figure 17). This loop is composed of several stations, including Denmark Hill and Elephant and Castle and deleting the proposed link (and thus breaking the loop) is the main factor that foster the trophic incoherence reduction.

5.2.2 Adding a link

Adding a link does not cause disruptions in the network, but it might not be possible due to physical / technical constraints. We measure the impact of a new link comparing its length with the average length of the links in the network. As a rule of the thumb, we assume that if a link is too long and there are not preexistent tracks, then the modification in the real network could be difficult or not convenient and we discard the option.

After simulating the insertion of each possible link, our analysis shows that adding a link from Bickley to Nunhead would be an optimal choice because it provides a significant decrease of the trophic incoherence. Adding
Figure 17: The highlighted loop (figure on the left, before the link removal) is responsible of the elevated trophic incoherence of the network. Removing the link highlighted in red reduces the trophic incoherence and reduces the trophic level of the stations (figure on the right, after the link removal).

This link improves the trophic incoherence of the network by 30.9%. It is longer than the link average (2.9x average length) but we believe that adding this particular link is possible since a track from Bickley and Nunhead is already available and thus it would just require to add a direct train from those two stations.

Figure 18: Stations trophic level. On the left the trophic level of the Thameslink stations, on the right the trophic level of the same stations when the link from Bickley to Nunhead is added.
**Trophic Coherence Analysis**  In Figure 18 we show the incoherence gain of the Thameslink network when a direct train service from Bickley to Nunhead is added. Denmark Hill, the node with the highest trophic level, decreases its level from 22.8 to 11 and West Bulwich becomes the node with the new highest trophic level (from 19 to 15). Obviously, adding a link does not remove any loop, but instead it reduces the trophic incoherence adopting data-driven (using morning peak-hour commutation data) changes in the network structure. In particular, this modification suggests a better design for the morning peak hours rides improving the connection between basal nodes (origin of the commuters) and top nodes (destination of the commuters).

**Data, Code and Materials**

Data available from the Dryad Digital Repository [32]:
Dryad review URL: [https://datadryad.org/review?doi=doi:10.5061/dryad.6s76rp7](https://datadryad.org/review?doi=doi:10.5061/dryad.6s76rp7)
Dryad DOI: [https://doi.org/10.5061/dryad.6s76rp7](https://doi.org/10.5061/dryad.6s76rp7)

**Competing Interests**

We have no competing interests.

**Authors’ Contributions**

W.G., L.V., S.J., & A.W. planned the experiments. A.P. & G.M. conducted the main numerical analysis. A.A. sourced the data and assisted in the analysis. A.P. & W.G. wrote the paper. All authors helped to analyze the findings and proof read the paper.

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